SCHEDULING OF SUCCESSIVE INTERFERENCE CANCELLATION (SIC) BASED WIRELESS AD HOC NETWORKS USING NAKAGAMI FADING CHANNEL

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ABSTRACT-Interference cancellation, congestion and better scheduled length these are the major issue of the wireless ad hoc network. In this research paper we have proposed a technique least interference effect-two ray ground (nakagami)(LIE-TRG(nakagami)) which is the modification of LIE technique. We have used nakagami fading channel to get better scheduled length with constant fading speed and two ray ground model to avoid the congestion in the network. SIC receiver is used for the better reception signals at the receiver side. SINR is used to calculate the signal to interference. We have used matlab 2014Rs to simulate our proposed work and the results are compared with the existing results.

Keywords: SIC, Nakagami, Two ray ground, LIE.

INTRODUCTION

This work offers an integrated outline to study the presentation of successive interference cancellation (SIC) in wireless networks with random fading circulation and power-law path loss. An investigative classification of the performance of SIC is assumed as a function of dissimilar scheme parameters. The results propose that the minimal advantage of permitting the receiver to successively decode users reduces very fast with, specifically in networks of high sizes and small path loss booster. On the other hand, SIC is extremely useful when the users are collected about the receiver and/or very low-rate codes are used. Similarly, with multiple packet reception, a lower per-user info rate constantly results in higher collective throughput in interference-limited networks.

In difference, there occurs a positive optimal per-user rate that exploits the aggregate throughput in noisy networks. The investigative results assist as useful tools to appreciate the possible gain of SIC in heterogeneous cellular networks (HCNs). Expending these implements, this work quantifies the improvement of SIC on the analysis likelihood in HCNs with non-accessible base stations. A stimulating reflection is that, for modern narrow-band schemes (e.g., LTE and WiFi), maximum of the gain of SIC is accomplished by canceling a single interferer.

PROBLEM STATEMENT

The problem is in the downlink providing the motivation to examine successive interference cancellation (SIC) as an interfering organization device LTE system and beyond. The Model exposed that well-ordered successive interference (OSIC) out achieves non-ordered successive interference cancellation (NSIC) and the added difficulty is acceptable built on the connected improvement in BER performance of OSIC. The main disadvantage of OSIC is that it is not effectual in network situation using power control or power provision. Other interference management methods will be compulsory to totally achieve the interfering. In [2], Mollanoori and Ghaderi deliberate the uplink scheduling problems for networks associate SIC, and resulting the optimum decoding instruction of the simultaneous transmissions. In this paper, we present an investigative framework for revising the presentation of SIC in large-scale-enabled cellular networks using the implements from stochastic geometry. To simplify the interference investigation, we recommend the method of stochastic correspondence of the interference, which changes the interference (interference from both the cellular


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tier and scheduling) to corresponding single-tier interference.

III. SYSTEM MODEL

We deliberate an ad hoc network with n nodes X1, X2, . . . , Xn. All nodes have a transmitter, a receiver, & an infinite buffer, and needs to connect with certain or entirely of the other nodes, perhaps by multi-hop routing. We undertake that nodes cannot communicate and receive at the same time, and that broadcasts reside in the full bandwidth W of the channel.

Node Xi can communicate with any power Pi ≤ Pmax. When Xi transmits, Xj obtains the signal with power GijPi, where Gij signifies the channel gain between nodes Xi and Xj. We describe the channel gain matrix to be the n × n matrix G = [Gij]. The origins beside the diagonal are insignificant, and are set to Gii = 0. We accept that the basics of the gain matrix are frequency-independent, but may be time-varying. The receiver of every node is issue to thermal noise, circumstantial interference from numerous noise sources such as additional networks, and interference from other users, where the interference initiated by Xi to Xj is also resolute by the link gain Gij. We perfect thermal noise and circumstantial interference equally, as a single source of noise, with power spectral density 𝜈i for node Xi. We describe the noise vector H = [η1 η2 . . . ηn]T.

Let {Xt : t ∈ J} be the set of communicating nodes at a assumed time, each node Xt communicating with power Pt. Let us take responsibility that node Xj, j ∈ J is trying to receive information from node Xi, i ∈ J. Then the signal to interference and noise ratio (SINR) at node Xj will be

\[ γ_{ij} = \frac{g_{ij} P_i}{\sum_{k \neq j} g_{ik} P_k + η_j} \]  (1)

We adopt that the rate of transmission from node Xi to node Xj, under the SINR γij given by (1), is given by Rij = fR(γij), where the receiver function fR(·) reflects the capabilities of the receiver and the performance metric. Many choices are possible for fR(·). For example, we can set

\[ f_R(η_j) = W \log_2 (1 + \frac{η_j}{T}) \]  (2)

where \( \log_2(x) \) signifies the logarithm of x, in base 2. Equation (2) infers that the transmitter incessantly adjusts the transmission rate to the receiver SINR.

With \( \Gamma = 1 \), the receiver achieves Shannon’s capacity. With \( \Gamma > 1 \), (2,3) approximates the maximum data rate that meets a given BER requirement under a specific modulation and coding scheme such as coded SIC [4,9]. The precise value of \( \Gamma \) depends on the choice of coding and modulation parameters, and the BER requirement.

Alternatively, we can use a receiver function that is more common in studies of medium access control and wireless LANs. In particular, we can assume that nodes use for all transmissions a common transmission rate R, and that the receiving node Xj will be able to decode the signal from Xi, with a negligible probability of error, provided the signal to interference and noise ratio γij is constantly above a given threshold γT : γij ≥ γT. Otherwise, the signal is lost. In other words,

\[ fR(γij) = R × 1 \{ γ_{ij} ≥ γT \} \]

This is a somewhat pessimistic assumption, as typically there is a non-negligible probability that a packet that is below but close to the SINR threshold will be correctly decoded. However, we will use this assumption for clarity of exposition, and none of our results will be significantly affected if a more accurate model is used.

IV. PROPOSED IMPLEMENTATION

Nakagami fading channel by multipath scattering with moderately large time delay spreads spreads with different clusters of reproduced waves. Within a cluster, the phase of separate waves are arbitrary but the time delays are almost equal for all the wave, which consequences in Rayleigh distribution for signals in each cluster. The average time delay is supposed to differ between the clusters. The magnitude of received signal is characterized by Nakagami distribution.

Channel Capacity- is the measure of maximum information that can be communicated constantly over a communication channel or maximum possible transmission rate when the probability of error is almost zero. Claude Elwood Shannon established the following equation for the theoretical channel capacity is

\[ C = B \log_2 (1+SNR) \]  (3)

Where C is capacity of channel in terms of bits/second, B is Bandwidth in Hertz, SNR is Signal to Noise ratio i.e. ratio of signal power to associated noise power.

Steps involved in proposed work:

1. Nakagami system is combined with Two Ray Ground system to enhance the capacity of the work.
2. Calculating the Network reliability of the channels.

3. Then proposed algorithm is applied. Firstly Nakagami which covers a wide range of fading conditions; when m=1/2, it is a one-sided Gaussian distribution and when m=1 it is a Rayleigh distribution. The two ray ground reflection model considers both the direct path and a ground reflection path. It is shown that this model gives more accurate prediction at a long distance than the free space model.

4. Finally we get the enhanced results i.e. capacity of the system is increased.

Flowchart of the steps performed

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LIE
TWO RAY GROUND MODEL
NAKAGAMI FADING CHANNEL
BINARY DECISION MAKING
FIND THE ARBITRARY NODES
NOW FIX ALL THE TERMINALS FOR EACH CHANNEL
CALCULATE THE SMALLEST DEGREE
CALCULATE THE LOG SHADOWING METHOD
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V. RESULT

The above graph shows that as the path loss exponent increases the proposed schedule length decreases showing that there is less path loss during transmission and also shows that there is an improvement in capacity of Fading channel when the Two ray ground Nakagami fading solution is implemented to achieve capacity maximization is used to allocate different powers to the sub channels.

Fig 5.2: Schedule length using TRG-fading for different path loss exponents on network coverage area.

In above graph the network coverage area is getting improved with network reliability at initial stage to gaining better improvement when coverage...
increasing in response time as compare to existing method.

Fig 5.4: Comparison between LIE and LIE-TRG (Nakagami) for Runtime on No of Links

In the above graph it is shown that LIE-TRG (Nakagami) technique is more efficient than LIE by comparing their runtime.

In this research some information about the Nakagami system in addition, problems related to MANET systems such as process interactions are presented. Also different strategies such as two ray ground method, to reduce these control interactions are presented.

VI. CONCLUSION

Using a Nakagami-Two ray ground established framework, this paper examines the performance of SIC in LIE-TRG Nakagami fading networks with power law density functions. We show that the probability of sequentially decoding at least network users decays super exponentially with high-rate codes are used, and it decays mainly fast under small path loss exponent in high dimensional networks, which proposes the negligible gain of adding more SIC capability reduces very fast. On the other hand, SIC is exposed to be principally beneficial if very low-rate codes are used or the dynamic transmitters are clustered around the receiver.

REFERENCES


BIOGRAPHY

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