

Performance Diagnosis Using Bigtable Monitoring for Cloud Computing System

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Abstract— Reduplication services are allocated in individual physical nodes in the cloud. They can be amassed into various types of services, serving large amount of user requests. The existing diagnosis techniques for such distributed systems can't effectively solve. In fine-grained performance abnormality is still lacking. Production cloud systems should be completely unaided. CloudDiag systematically collects the point-to-point tracking data from each physical node in the cloud and classifies into different categories according to call tree. It then employs a customized Map-Reduce algorithm to proactively analyze the tracing data. Specifically, it assembles the tracing data of each user request, and classifies the tracing data into different categories according to call trees of the requests.

Index Terms— abnormality, BigTable, Cloud, CloudDiag Performance diagnosis, Pont-to-Point.

I. INTRODUCTION

Cloud computing system is most important thing in this world. We can use a service without having, which is a milestone in networking field. But, some time we are faced many problems in networking field. Specifically we face security problems in cloud computing system. When server serves the service to clients in cloud computing, some performance diagnosis problems are occurs. Server may not be fulfilling all user requests due to network traffic, unaided diagnosis tool for locating fine-grading performance abnormality entries are still lacking on cloud computing system. Performance diagnosis tool for production cloud systems should be completely unaided. This paper bridges this gap by proposing CloudDiag. CloudDiag periodically collects the point-to-point tracing data (In particular, execution time of method invocations) from each physical node in the cloud. When the cloud system is suffering performance degradation (e.g., average response time of user requests is larger than a threshold), a cloud operator can access CloudDiag with its web interfaces to conduct a performance diagnosis. With the request tracing data, CloudDiag will perform a fast customized matrix recovery algorithm to instantly identify the method invocations

(together with the reduplication they locate) which contribute

the most to the performance anomaly. The whole process requires no domain-specific knowledge to the target service.

A. Fine Grained Performance model for Cloud Computing

Performance diagnosis is labor intensive in production cloud computing systems. Commonly Such systems face many more real-world challenges, which existing diagnosis techniques for such distributed systems cannot effectively solve. Combining a statistical technique and a fast matrix recovery algorithm, CloudDiag can efficiently mistakes fine-grained causes of the execution problems, which does not require any domain-specific ability to the target system. CloudDiag has been utilized in a realistic production cloud computing systems to diagnose performance problems. We determine the capability of CloudDiag in three real-world case studies.

Accurate performance evaluation of cloud computing resources is a necessary prerequisite for ensuring that quality of service parameters remain within agreed limits. In this paper, we employ both the examining and simulation modeling to addresses the complexity of cloud computing systems. Examining model is comprised of distinct functional sub models, the results of which are combined in an iterative manner to obtain the solution with need accuracy. Our models incorporate the features of cloud centers such as resource virtualization, batch arrival of user requests, and realistic servicing steps, obtain important accomplishment metrics such as task blocking probability and total waiting time incurred on user requests. Also, our results reveal substantial awareness for capacity planning to control delay of servicing users' requests.

B. Performance Data Collection

What kind of performance data that CloudDiag should collect and how to collect them. Our instrumentation-based tracing approach will produce performance data when a sampled request is being processed in each component reduplication. Specifically, each component method, when being request or returning, will generate a log entry. Host indicates the machine where the component reduplication locates. Time stamp records the time of the event occurrence. Request ID is the global identifier of a request. MID is a unique identifier for request tracing purpose, which will be discussed later. The Method field saves the name of the

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method invoked. Lastly, Flag indicates whether this is a method invocation or a method return.

C. Performance Diagnosis in Fine Granularity

A component typically has a lot of reduplication in a production cloud system. Within one component, there are many performance-related private methods (those invoked inside the component) and public methods (the interfaces invoked by other components). It is very challenging to localize anomalous methods as well as their corresponding physical reduplication. Current approaches generally focus on locating anomalous physical nodes. Such coarse-grained results are not enough. In the former case, given an anomalous physical node, a system operator has to identify the faulty component among many components typically hosted in the same node. In the latter case, given an anomalous logical component, the operator has to identify which one among their numerous reduplication distributed in the cloud is faulty. Consequently, huge human efforts are still required to further pinpoint the subtle primary cause. Performance diagnosis in a fine granularity is of high concern to reduce manual efforts.

D. Unaided Performance Diagnosis

CloudDiag has been successfully launched in diagnosing performance problems for the production cloud systems in Cloud Computing. We report three case studies in our real-world performance diagnosis experiences to demonstrate the effectiveness of CloudDiag in helping the operators localize the primary causes of performance problems. Many existing performance diagnosis techniques resort to system behavior models in identifying abnormality. Unfortunately it is hard to manually build such models in production cloud systems, given their complexity in system scale. In addition, cloud services are generally composed of many components developed by different teams, which are independently updated online. It is extremely difficult to maintain the behavior models for such evolutionary systems. Hence, a performance diagnosis tool for production cloud systems should be completely unaided, without assuming that any prior knowledge about the service should be input.

The goal of cloud computing is to apply traditional supercomputing, or high-performance computing power, normally used by military and research facilities, to perform tens of trillions of computations per second, in consumer-oriented applications such as financial portfolios, to deliver personalized information, to provide data storage or to power large, immersive computer games. To do this, cloud computing uses networks of large groups of servers typically running low-cost consumer PC technology with specialized connections to spread data-processing chores across them. This shared IT infrastructure contains large pools of systems that are linked together. Often, virtualization techniques are used to maximize the power of cloud computing.

II. SURVEY ON PERFORMANCE DIAGNOSIS

A. Resource Monitoring, Low-Level Metrics Mapping, and Detection in Cloud Computing

Z. Lan., Z. Zheng Ref. [2] cloud computing represents a novel paradigm for the implementation of scalable computing infrastructures combining concepts from virtualization, distributed application design, Grid, and enterprise IT management. Service Level Agreements (SLAs) representing a contract signed between the customer and the service provider including non-functional requirements of the service specified as Quality of Service (QoS). There is little work in the area of resource monitoring, low-level metrics mapping, and SLA violation detection in Cloud computing. Because of that, we look into the related areas of Grid and Service-Oriented Architecture (SOA) based systems. FoSII project (Foundations of Self-governing Infrastructures) is developed models and concepts for autonomic SLA management and enforcement in Clouds.

In order to save Cloud providers from paying penalties and increase their profit, providers have to monitor the current status or resource and check frequently whether the established SLAs are violated. The main components of FoSII architecture are the automatic VM deployed, responsible for the allocation of resources and for mapping of tasks, application deployed, responsible for the execution of user applications, deployed a pinpointing system which represents performance deviations in subsystems. Their work uses performance counter data of a load test to craft performance signatures for the LSS subsystems which performs principal component analysis, to reduce the large volume of performance counter data. However, their technique cannot be generalized to other domains such as network traffic and security monitoring. This is due to the fact that there is no guarantee that the directions of maximum variance will contain good variables for discrimination. A large anomaly may inadvertently pollute the normal subspace, thereby skewing the assumption that large variances always have important dynamics.

One of the contribution in Online System Problem Detection by Mining Patterns of Console Logs is a novel two-stage online log processing approach that combines frequent pattern mining with Principal Component Analysis (PCA) based anomaly detection for system runtime problem detection. Frequent pattern mining and distribution estimation techniques to capture the dominant patterns have been used to detect components failure. This online logs used to identify and filter out common (normal) operational patterns from free-text console logs, and then perform PCA based anomaly detection on the remaining patterns to identify operational problems within minutes of their occurrence (as represented by information in the console logs).

Google has developed several analysis tools for use with Dapper. Most relevant is the Service Inspector, which shows graphs of the unique call paths observed to a chosen function

or component, along with the resulting call tree below it, allowing developers to understand the contexts in which the chosen item is used. Comparing request flows, as captured by end-to-end traces, is a powerful new technique for diagnosing performance changes between two time periods or system versions. Spectroscope's algorithms have been proposed by Alice X. Zheng for this comparison, allow it to accurately identify and rank mutations and identify their precursors, focusing attention on the most important differences.

B. Detecting Large-Scale System Problems by Mining Console Logs

Wei Xu, Ling Huang Dec 2012 Ref [3], Tracing real requests through the system enables us to support problem determination in dynamic systems where using dependency models is not possible. This tracing also allows us to distinguish between multiple instances of what would be a single logical component in a dependency model. By performing data clustering to analyze the successes and failures of requests, Mike Y. Chen, Emre Kiciman, Eugene Fratkin, have attempted to find the combinations of components that are most highly correlated with the failures of requests, under the belief that these components are causing the failures. By analyzing the components that are used in the failed requests, but are not used in successful requests, they provide high accuracy with relatively low number of false positives. This analysis detects individual faulty components, as well as faults occurring due to interactions among multiple components.

This framework, Pinpoint, requires no application-level knowledge of the systems being monitored or any knowledge of the requests. This makes Pinpoint suitable for use in large and dynamic systems where this application-level knowledge is difficult to accurately assemble and keep current. As such, it is an important improvement over existing fault management approaches that require extensive knowledge about the systems being monitored. Pinpoint traces requests as they travel through a system, detects component failures internally and end-to-end failures externally, and performs data clustering analysis over a large number of requests to determine the combinations of components that are likely to be the cause of failures. The runtime tracing and analysis is necessary for systems that are large and dynamic, such as today's Internet systems.

The accuracy in log parsing allows us to extract the identifiers and state variables, which are widely found in logs yet are usually ignored due to difficulties in log parsing. Using console logs, we are able to construct powerful features that were previously exploited only in structured traces. Detecting Large-Scale System Problems by Mining Console Logs has opened up many new opportunities for turning built-in console logs into a powerful monitoring system for problem detection, and suggests a variety of future directions that can be explored, including: 1) extracting log templates from program binaries instead of source code, which not only

makes our approach work on non-open-source modules but also brings much operational convenience; 2) designing other features to fully utilize the rich information in console logs; 3) developing online detection algorithms instead of current *postmortem* analysis; and 4) investigating methods to correlate logs from multiple related applications and detect more complex failure cases.

C. End-to-End Request-Flow Tracing

M. Chen, E. Kiciman Nov 2010 Ref [13], Request-flow comparison builds on end-to-end tracing, an invaluable information source that captures a distributed system's performance and control flow in detail. Such tracing works by capturing activity records at each of various trace points within the distributed system's software, with each record identifying the specific trace point name, the current time, and other contextual information. Most implementations associate activity records with individual requests by propagating a per-request identifier, which is stored within the record. Activity records can be stitched together, either offline or online, to yield request-flow graphs, which show the control flow of individual requests. Several efforts, including Magpie, Whodunit, Pinpoint, X-Trace, Google's Dapper, and Stardust have independently implemented such tracing and shown that it can be used continuously with low overhead, especially when request sampling is supported. For example, Stardust, Ursa Minor's end-to-end tracing mechanism, adds 1% or less overhead when used with key benchmarks, such as SpecSFS.

D. Spectroscope

To illustrate the utility of comparing request flows, this technique was implemented in a tool called Spectroscope and used to diagnose performance problems seen in Ursa Minor [1] and in certain Google services. This section provides an overview of Spectroscope, and the next describes its algorithms. Even small distributed systems can service hundreds to thousands of requests per second, so comparing all of them individually is not feasible. Instead, exploiting a general expectation that requests that take the same path should incur similar costs, Spectroscope groups identically-structured requests into unique categories and uses them as the basic unit for comparing request flows.

For example, requests whose structures are identical because they hit in a NFS server's data and metadata cache will be grouped into the same category, whereas requests that miss in both will be grouped in a different one. Two requests are deemed structurally identical if their string representations, as determined by a depth first traversal, are identical. For requests with parallel substructures, Spectroscope computes all possible string representations when determining the category in which to bin them. The exponential cost is mitigated by imposing an order on parallel substructures (i.e., by always traversing them in alphabetical order, as determined by their root node names) and by the fact that parallelism is limited in most request flows we have observed.

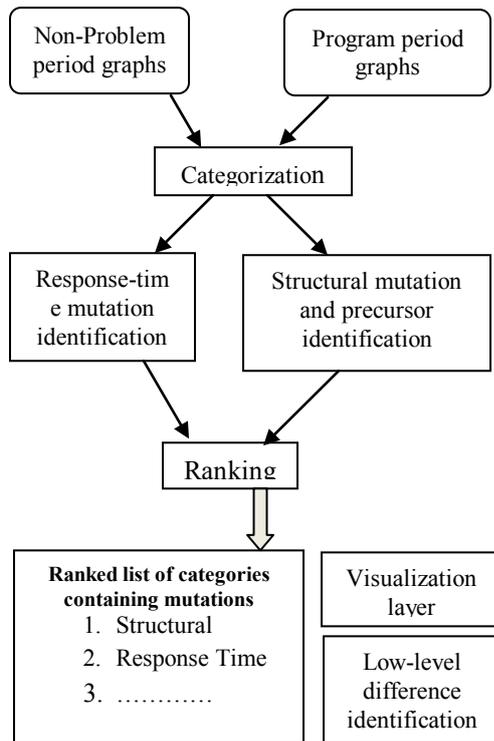


Figure 2.1 Spectroscope's workflow

First, Spectroscope groups requests from both periods into categories. Second, it identifies which categories contain mutations or precursors. Third, it ranks mutation categories according to their expected contribution to the performance change. Developers are presented this ranked list. Visualizations of mutations and their precursors can be shown. Also, low-level differences can be identified for them. For example, in Ursa Minor, one performance slowdown, which manifested as many structural mutations, was caused by a change in a parameter sent by the client. For problems like this, highlighting the specific low-level differences can immediately identify the root cause. Section 5 describes Spectroscope's algorithms and heuristics for identifying mutations, their corresponding precursors, their rank based on their relative influence on the overall performance change, and their most relevant low-level parameter differences. It also describes how these methods overcome key challenges—for example, differentiating true mutations from natural variance in request structure and timings. Identification of response-time mutations and ranking rely on the expectation (reasonable for many distributed systems, including distributed storage) that requests that take the same path through a distributed system will exhibit similar response times and edge latencies.

E. Distributed Storage System for Structured Data

F.Chang, J.Dean, Aug 2012 Ref [12], Bigtable is a distributed storage system for managing structured data that is designed to scale to a very large size: petabytes of data across thousands of commodity servers. Many projects at Google store data in Bigtable, including web indexing, Google Earth, and Google Finance. These applications place very different demands on Bigtable, both in terms of data size

(from URLs to web pages to satellite imagery) and latency requirements (from backend bulk processing to real-time data serving). Despite these varied demands, Bigtable has successfully provided a high-performance solution for all of these Google products. In this paper we describe the simple data model provided by Bigtable, which gives clients dynamic control over data layout and format, and we describe the design and implementation of Bigtable.

Over the last two and a half years we have designed, implemented, and deployed a distributed storage system for managing structured data at Google called Bigtable. Bigtable is designed to reliably scale to petabytes of data and thousands of machines. Bigtable has achieved several goals: wide applicability, scalability, high performance, and high availability. Bigtable is used by more than sixty Google products and projects, including Google Analytics, Google Finance, Orkut, Personalized Search and Google Earth. These products use Bigtable for a variety of demanding workloads, which range from throughput-oriented batch-processing jobs to latency-sensitive serving of data to end users.

The Bigtable clusters used by these products span a wide range of conjurations, from a handful to thousands of servers, and store up to several hundred terabytes of data. In many ways, Bigtable resembles a database: it shares many implementation strategies with databases. Parallel databases and main-memory databases have achieved scalability and high performance, but Bigtable provides a different interface than such systems. Bigtable does not support a full relational data model; instead, it provides clients with a simple data model that supports dynamic control over data layout and format, and allows clients to reason about the locality properties of the data represented in the underlying storage. Data is indexed using row and column names that can be arbitrary strings. Bigtable also treats data as uninterrupted strings, although clients often serialize various forms of structured and semi-structured data into these strings. Clients can control the locality of their data through careful choices in their schemas. Finally, Bigtable schema parameters let clients dynamically control whether to serve data out of memory or from disk.

III. PROBLEM DEFINITION

A system operator has to identify the faulty component among many components typically hosted in the same node. The operator has to identify which one among their numerous reduplication distributed in the clo0075d is faulty. Huge human efforts are still required to further pinpoint the subtle primary cause. Unaided performance diagnosis is hard to manually build such models in production cloud systems, given their complexity in system scale. A component typically has a lot of reduplication in a production cloud system. However, the performance anomaly of a component may be manifested only in a small part of its reduplication. This will cause the performance degradation of the involving

service, which is frequently observed in the cloud computing systems.

IV. MODULE DISCRPTION

An efficient, unaided diagnosing tool that can pinpoint the fine-grained causes of performance abnormality is of critical importance to production cloud computing systems. To this end, we propose CloudDiag, a tool for performance diagnosis in production cloud computing systems. CloudDiag is composed of three major parts, i.e., 1) collecting the performance data; 2) assembling the performance data; and 3) identifying the primary causes of the abnormality. CloudDiag conducts the tracing data collection and assembly during the execution of the system. All categories are stored in a Big Table-like storage system for further performance diagnosis when performance abnormalities are detected.

A. Reduplication Implementation

A service with abnormality performance must have involved some components with performance abnormality. A component typically has a lot of reduplication in a production cloud system; however, the performance anomaly of a component may be manifested only in a small part of its reduplication. Data replication if the same data is stored on multiple storage devices. Computation replication if the same computing task is executed many times. If we replicate data, these processes are passive and operate only to maintain the stored data, reply to read requests, and apply updates. When we replicate computation, the usual goal is to provide fault-tolerance.

B. Users Request Trace

Performance problems are the most difficult to locate, because the anomalous methods hide themselves in numerous well-functioning reduplication. Therefore, to reduce human efforts in pinpointing performance anomaly, a performance diagnosis tool must first identify which component methods contribute to the performance anomaly, and then locate the component reduplication that execute the methods. CloudDiag conducts the tracing data collection and assembly during the execution of the system. CloudDiag traces user requests at a given sampling rate to expose performance data. For the sampled requests, each component replica records the performance data and saves them in its local storage. An important consideration is what kind of performance data. CloudDiag adopts an instrumentation-based approach that collects the execution time of each component method.

C. Bigtable-Like Storage System

CloudDiag conducts the tracing data collection and assembly during the execution of the system. All categories are stored in a Big-Table-like storage system for further performance diagnosis when performance abnormalities are detected. A web-based interface to access the data is provided for system operators. When performance abnormalities are observed, an operator can conduct the primary cause analysis by accessing the web interface. Data is indexed using row and column names that can be arbitrary strings. Big-table also

treats data as uninterested strings, although clients often serialize various forms of structured and semi-structured data into these strings. Clients can control the locality of their data through careful choices in their schemas. Finally, Bigtable schema parameters let clients dynamically control whether to serve data out of memory or from disk.

D. Abnormality Identification

To mark the distributed tracing record and use their time stamps to infer the invocation relationships with methods. As a result, clock drifts may lead to the wrong order of method invocations when merging the distributed logs generated in processing a request. To guarantee the order of the invoked methods, we design hierarchical identifiers to trace a request. Specifically, before a node calls a method in another node, besides passing the Request ID to the caller, the caller also generates an MID, a unique integer, for the caller. When a request enters the first component method in the system (i.e., the request entry method), the MID of the method is initially set identical to the Request ID. CloudDiag then records the MIDs of the caller and the caller in the logs of the caller as well. A customized map-reduce process is utilized to group requests into different categories based on their call trees. Requests within one category share the same call tree. CloudDiag then identifies the anomalous categories according to their latency distribution. Then, for each anomalous category, a fast customized matrix recovery algorithm is employed to identify the anomalous method invocations together with the reduplication they are located.

E. CloudDiag Architecture

The following systems are there in CloudDiag Architecture Fig. 4.1. That is consumers, cloud tracking host, log, public cloud, and big table storage system. The consumers are given a request to the cloud tracking host for public cloud service. In that time cloud tracking host give a response to the consumer that is nothing but, give a permission to access the service from the public cloud. Reduplication services are used to serve more than one user at the same time. That data will be stored in the big table storage system and other maintained in a log. Big table storage is used to maintain a users and finding the abnormality entries in the system. Here we give a time limit for a each service that is helped to omit such users, who is not responding to the cloud tracking host. Each user has a unique resource so we can identify user and service related problems easily.

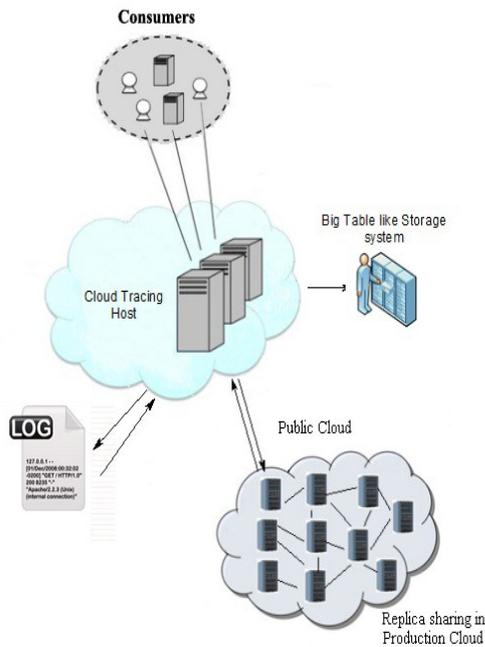


Figure 4.1 CloudDiag Architecture

F. Sequence of the cloudDiag

Sequence of the cloudDiag is a complete working principle of the cloudDiag. First of all the servers receive a request from the consumer, then find configuration of the client system. Because some cloud services are no platform independent, user may use various operating systems like Windows, Linux, MacOS, and etc. Each user have a unique task, server allocate resource for each task. Computing a task of resources. Some task may be giving a failure reports. So, server monitors a resource failure. Then the sequence is divided into two parts. The first one is the server responding the process task then give a response to corresponding user or client. The other one is server rescheduling the reduplication services and compute the reduplication service and filled in the log entry. Then give a response to corresponding user or client as usual. Finally the server identifies the abnormal methods in host and stored in a bigTable storage system. User information, resource information and service information are stored and maintained in a database like bigTable.

V. CONCLUSION

CloudDiag resorts to a white-box instrumentation mechanism to trace service requests, because the source codes of services are generally available in typical production cloud systems. Note that such a white-box performance data acquisition component of CloudDiag can also be substituted with another tracing mechanism if it can obtain the latency data of method invocations. There is a tradeoff between tracing granularity and debugging effort. As a result, more effort will be increased in troubleshooting the performance abnormality if a black-box tracing mechanism is applied.

Our future concentration is to explore black-box tracing mechanisms so that a fine granularity, which incorporate black-box tracing mechanisms with CloudDiag. To this end,

the runtime instrumentation can be a promising technique. In the future, we also to extend it with the capability to manage a Cloud environment with multiple data centers. Thus, we will apply a decentralization approach whereby the proposed system is installed on each data center.

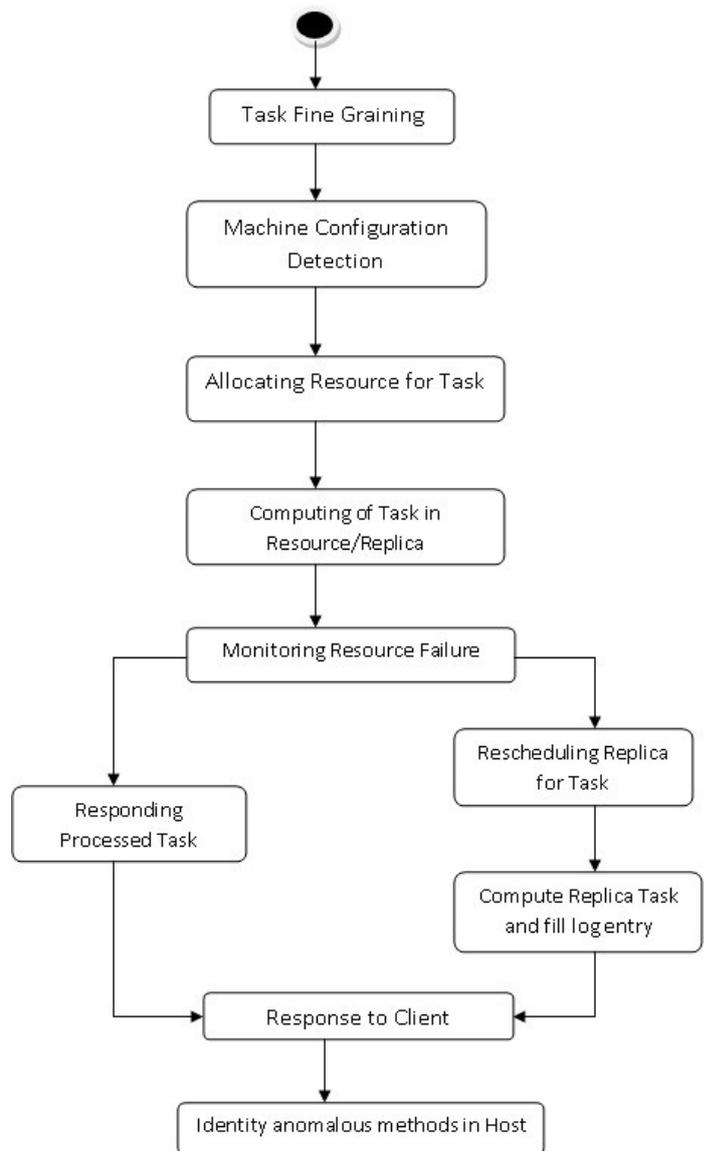


Figure 4.2 CloudDiag Sequence

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