

Overview of Sequential Pattern-Mining Algorithms

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Abstract- The concept of Sequential Pattern Mining was first introduced by R. Agrawal and R. Srikant (1995). It aimed to retrieve frequent patterns in the sequences of products bought by customers through time ordered transactions. Sequential pattern mining is significant data-mining method for determining time-related behaviour in sequence databases. In this paper, a survey of the sequential pattern mining algorithms is performed. This paper represents a review on various algorithms of discovering sequential patterns from large sequence databases which is an important problem in the field of knowledge discovery and data mining.

Keywords: - sequential patterns, data mining, knowledge discovery, sequential pattern mining.

I. INTRODUCTION

The Sequential pattern mining was first proposed by Agrawal and Srikant [1] and later many algorithms were proposed by many others for performance improvements. Initially the inventors refined their algorithm by introducing Generalized Sequential Patterns (GSP) based on the based on implementing apriori dimensions [2]. Few other interesting algorithms on persistent pattern mining are SPADE [3], SPAM [4], and PrefixSpan [5]. The SPADE uses lattice theoretic approach based on vertical ID list in order to optimize original search space into optimized smaller spaces. Later SPAM was introduces for mining long patterns this algorithm espouses a vertical bitmap representation and the performance scale shows that SPADE is more efficient than SPAM but still there is a flaw in this algorithm as it consumes more space compared with the SPAM algorithm.

II. APRIORI-BASED METHODS

A. AprioriAll

AprioriAll [1] scans the database several times to find frequent itemsets of size k at each k^{th} -iteration (starting from $k = 2$). It also has the generate-and-test feature by performing the *Apriori-generate* join procedure to join L_{k-1} with itself to generate C_k , the set of candidate sequences in the k^{th} -iteration, it then prunes sequences in C_k which have subsequences not in L_{k-1} (i.e., are not large), creates L_k by adding all sequences from C_k with support $\geq \text{min sup}$ until

there are no more candidate sequences. The pruning procedure performed on itemsets in C_k removes sequences that have infrequent subsequences. The AprioriAll technique suffers from increased delays in mining as the number of sequences in the database gets larger.

B. PSP:

This is another apriori-based algorithm, also built around GSP, but the difference of using a prefix-tree. In PSP [6], user access sequences are sorted in the web log according to the IP address. The user is allowed to provide a time period t by which access sequences that are temporally close to each other are grouped. A prefix-tree is then built to handle the mining procedure in a way similar to GSP. Following up on tree traversal, graph traversal mining uses a simple unweighted graph to reflect the relationship between pages of websites. The algorithm is similar to apriori, without performing the Apriori-generate join. The database still has to be scanned several times, but it is more efficient than GSP.

C. GSP:

The GSP algorithm, adopt a multiple-pass candidate generate-and-test method for finding sequential patterns. The first pass on the database determines the support for each item in finding frequent 1-sequences based on the minimum support. The first ($k = 1$) scan over the table generates the set of candidate 1-sequences C_1 , giving the seed set of frequent 1-sequences L_1 . These sequences provide a *seed set* L_k which is used to generate next-level candidate sequences in the next pass $k+1$. The next C_{k+1} candidate set is generated by performing a GSP-join of L_k on itself. The GSP-join, like the apriori-generate join requires that two sequences in L_k join together if two conditions are met. Here, a sequence s_1 in L_k joins another sequence s_2 in L_k if the subsequence obtained by dropping the first $(k - 1)$ items of s_1 is the same as the subsequence obtained by dropping the last $(k-1)$ items of s_2 , hence the candidate sequence generated is the sequence s_1 extended with the last item in s_2 and is added to C_k . The prune phase deletes candidate sequences that have a contiguous $(k - 1)$ -subsequence with support less than *min sup*. Each candidate sequence has one more item than a seed k -sequence, so all candidate sequences have the same number of items at each level. Support for these candidate sequences is again found during a pass over the data. The algorithm terminates when

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no new sequential pattern is found in a pass, or no candidate sequence can be generated. The generate-and-test feature of candidate sequences in GSP consists of two phases: Join Phase and Prune Phase. During the join phase, candidate sequences are generated by joining L_{k-1} with itself using *GSP-join*. For increased efficiency, GSP employs a hash-tree to reduce the number of candidates in C that are checked for sequences. GSP is reported to be 2 to 20 times faster than AprioriAll.

D. SPAM:

SPAM is provided in Ayers et al. [4] integrates the ideas of GSP, SPADE, and FreeSpan. The entire algorithm with its data structures fits in main memory, and is claimed to be the first strategy for mining sequential patterns to traverse the lexicographical sequence tree in depth-first fashion. Each node in the tree has sequence-extended children sequences generated in the *S-Step* of the algorithm, and itemset-extended children sequences generated by the *I-Step* of the algorithm at each node. SPAM traverses the sequence tree in depth-first search manner and checks the support of each sequence-extended or itemset-extended child against *min sup* recursively. If the support of a certain child s is less than *min sup*, there is no need to repeat depth-first search on s by the apriori property. Apriori-based pruning is also applied at each S-Step and I-Step of the algorithm, minimizing the number of children nodes and making sure that all nodes corresponding to frequent sequences are visited. For efficient support-counting, SPAM uses a vertical bitmap data structure representation of the database similar to the id list in SPADE. Each bitmap has a bit corresponding to each element of the sequences in the database. Each bitmap partition of a sequence to be extended in the S-Step is first transformed using a lookup table, such that all the bits after the index of the first "1" bit (call it index y) are set to one and all the bits with index less than or equal to y are set to zero. When SPAM was compared to SPADE, it was found to outperform SPADE by a factor of 2.5, while SPADE is 5 to 20 times more space-efficient than SPAM, making the choice between the two a matter of a space-time trade-off. Here, Ayres et al. [4] propose to use a compressed bitmap representation in SPAM to save space, but do not elaborate on the idea.

III. PATTERN GROWTH BASED METHODS

A. FreeSpan:

Frequent pattern projected Sequential pattern mining [7] uses frequent items to recursively project sequence databases into a set of smaller projected databases and grows subsequence fragments in each projected database. This process partitions both the data and the set of frequent patterns to be tested, and confines each test being conducted to the corresponding smaller projected database. FreeSpan first scans the database, collects the support for each item, and finds the set of frequent items. Frequent items are listed in support descending order (in the form of item :support). According to flist, the complete set of sequential patterns in S can be divided into 6 disjoint subsets: (1) the ones containing only item 'a', (2) the ones containing item 'b', but containing no items after 'b' in flist, (3) the ones

containing item 'c', but no items after 'c', in flist, and so on, and finally, (6) ones containing item 'f'. The subsets of sequential patterns can be mined by constructing projected databases. Infrequent items, such as 'g' in this example, are removed from construction of projected databases.

B. PrefixSpan:

Prefix-projected Sequential pattern mining [5] works similar to FreeSpan except that the partitioning is done using prefixes of sequences. Its general idea is to examine only the frequent prefix subsequence's and project only their corresponding postfix subsequence's into projected databases because any frequent subsequence can always be found by growing a frequent prefix. No candidate sequence needs to be generated by PrefixSpan. Projected databases keep shrinking. The major cost of PrefixSpan is the construction of projected databases. To further improve mining efficiency, two kinds of database projections are explored: level-by-level projection and bi-level projection. Moreover, a main-memory-based pseudo-projection (using pointers rather than physically copying postfix sequences) technique is developed for saving the cost of projection and speeding up processing when the projected (sub)-database and its associated pseudo-projection processing structure can fit in main memory. PrefixSpan mines complete set of patterns much faster than both GSP and FreeSpan.

C. WAP-mine:

It is Pattern-Growth Miner with Tree Projection. At the same time as FreeSpan and PrefixSpan in 2000/2001, another major contribution was made as a pattern growth and tree structure-mining technique, that is, the WAP-mine algorithm [7] with its WAP-tree structure. Here the sequence database is scanned only twice to build the WAP-tree from frequent sequences along with their support; a "header table" is maintained to point at the first occurrence for each item in a frequent itemset, which is later tracked in a threaded way to mine the tree for frequent sequences, building on the suffix. The first scan of the database finds frequent 1-sequences and the second scan builds the WAP-tree with only frequent subsequences. The frequent subsequence of each original database sequence is created by removing infrequent 1-sequences. The tree starts with an empty root node and builds downward by inserting each frequent subsequence as a branch from root to leaf. During the construction of the tree, a header link table is also constructed, which contains frequent 1-sequences, each one with a link connecting to its first occurrence in the tree and threading its way through the branches to each subsequent occurrence of its node type. To mine the tree, WAP-mine algorithm starts with the least frequent item in the header link table and finds frequent items building on the suffix to get frequent k -sequences. The WAP-mine algorithm is reported to have better scalability than GSP and to outperform it by a margin. Although it scans the database only twice and can avoid the problem of generating explosive candidates as in apriori-based and candidate generate-and-test methods, WAP-mine suffers from a memory consumption problem, as it recursively reconstructs numerous intermediate WAP-trees during mining, and in particular, as the number of mined frequent patterns increases. This problem was solved by the PLWAP

algorithm [3], which builds on the prefix using position-coded nodes.

D. FS-Miner:

FS-Miner is a tree projection pattern growth algorithm that resembles WAP-mine and supports incremental and interactive mining [8]. The significance of FS-Miner is that it starts mining immediately with 2-subsequences from the second (which is also the last scan) of the database (at $k = 2$). It is able to do so due to the compressed representation in the FS-tree, which utilizes a header table of edges rather than single nodes and items, compared to WAP-tree and PLWAP-tree. It is also considered a variation of a trie, as it stores support count in nodes as well as edges of the tree that represents 2-sequences and is required for the incremental mining process. The sequence database is scanned once to find counts for links and to insert them in the header table. The FS-tree is built during the second scan of the database in a manner similar to WAP-tree and PLWAP-tree, except that a sequence that contains infrequent links is split and each part is inserted in the tree separately. Pointer links from header table are built to connect occurrences of header table entries, as they are inserted into the FS-tree similar to WAP-tree linkage.

IV. CONCLUSION

This survey has provided an overview of recent work in the area of sequential pattern mining algorithm techniques. Sequential pattern mining is trying to find the relationships between occurrences of sequential events, to find if there exist any specific orders of the occurrences. We can find the sequential patterns of specific individual items; also we can find the sequential patterns cross different items. With over a decade of extensive research, there have been hundreds of research publications and tremendous research, development and application activities in this domain.

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